

Adding structure to agent programming languages

(programming agents with mental states)

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Agenda

1 Motivation

- Agent programming languages
- 2 Structuring a rule-based language
 - Core language
 - Adding structure

3 Conclusion

- Notes and related work
- Ongoing & future work
- Contributions



Problem

Ideal agent oriented programming language:

modularity

- knowledge representation
- source code
- code encapsulation
- decomposition and combination ~~> compound structures
- code reuse!
- readability



The way to go...

1 basis for the minimalistic core language

- Modular BDI Architecture [Novák and Dix, 2006]
- IMPACT [Subrahmanian et. al., 2000]
- **2** structural extension:
 - alternative semantical view on the language
 - notion of mental state transformer
- **3** further extensions: *macros*



Modular core language: syntax

Definitions

Let ${\mathcal L}$ be a language:

- mental state is a theory $\sigma \subseteq \mathcal{L}$,
- query language \mathcal{L}_Q : $\varphi \in \mathcal{L}$, then $Q(\varphi) \in \mathcal{L}_Q$, $\top, \bot \in \mathcal{L}_Q$ and \lor, \land, \neg are allowed in \mathcal{L}_Q ,
- update language \mathcal{L}_U : $\varphi \in \mathcal{L}$, then $U(\varphi) \in \mathcal{L}_U$,
- let $\phi \in \mathcal{L}_Q$, $\psi \in \mathcal{L}_U$, then $\phi \longrightarrow \psi$ is a transition rule.

program ~~> a set of transition rules



Core language: standard operational semantics

Definitions

abstract operators:

$$Query_{\mathcal{L}} : 2^{\mathcal{L}} \times \mathcal{L} \to \{ true, false \}, (\sigma \models \varphi)$$

•
$$Update_{\mathcal{L}}: 2^{\mathcal{L}} \times \mathcal{L} \to 2^{\mathcal{L}}$$
, $(\sigma \oplus \psi = \sigma')$

Definition

Application of a rule $\phi \longrightarrow \psi$ in a state σ :

$$\frac{\sigma \models \phi, \sigma \oplus \psi = \sigma'}{\sigma \longrightarrow \sigma'}$$

Definition

Agent system semantics \rightsquigarrow set of all possible *computation runs* $\sigma_0, \sigma_1 \dots$ induced by the program \mathcal{P} .

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Core language: denotational semantics

Definition

Program \mathcal{P} can be characterized by a partial function

$$\mathcal{F}_{\mathcal{P}}: 2^{\mathcal{L}} \times \mathcal{L}_U \longrightarrow 2^{\mathcal{L}}; \mathcal{F}_{\mathcal{P}}(\sigma, \psi) = \sigma'$$

- σ is such, that \exists a rule $r = (\phi \longrightarrow \psi) \in \mathcal{P}$ and r is applicable in σ
- a corresponding set of states $\Sigma_{\mathcal{F}_{\mathcal{P}}} \subseteq 2^{2^{\mathcal{L}}}$ is an *application* domain of $\mathcal{F}_{\mathcal{P}}$

Program \mathcal{P} :

operational semantics: set of all the specified paths in the space of all possible mental states denotational semantics: set of all the specified transitions between classes of mental states

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Mental state transformer

Definition

A mental state transformer τ characterized by a function \mathcal{F} :

1 primitive mst: $\tau = \{\phi \longrightarrow \psi\} \rightsquigarrow \mathcal{F}(\sigma, \psi) = Update_{\mathcal{L}}(\psi, \sigma)$ and $\Sigma_{\mathcal{F}} = \{\sigma | \sigma \models \phi\}$

2 specialization: $\tau = \{\phi \longrightarrow \tau'\} \rightsquigarrow \mathcal{F}(\sigma, \psi) = \mathcal{F}'(\sigma, \psi)$ and $\Sigma_{\mathcal{F}} = \{\sigma | \sigma \in \Sigma_{\mathcal{F}'} \land \sigma \models \phi\}$

3 generalization: $\tau = \tau' \cup \tau'' \rightsquigarrow$ $\mathcal{F}(\sigma, \psi) = \begin{cases} \mathcal{F}'(\sigma, \psi) & \text{if } \sigma \in \Sigma_{\mathcal{F}'} \\ \mathcal{F}''(\sigma, \psi) & \text{if } \sigma \in \Sigma_{\mathcal{F}''} \end{cases} \text{ and } \Sigma_{\mathcal{F}} = \Sigma_{\mathcal{F}'} \cup \Sigma_{\mathcal{F}''}$



Example

when [{at(X,Y)}] and not [{desiredPosition(X,Y)}] and [{towards((X1,Y1),(X,Y))}] and not [{obstacle(X1,Y1)}] then [{stepTo(X1,Y1)}]

when [{at(X,Y)}] and [{goldAt(X,Y)}] and [{loaded}]
then [{broadcast(goldAt(X,Y))}]

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Example

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then [{broadcast(goldAt(X,Y))}]



define movementToPosition {

```
when [{at(X,Y)}] and not [{desiredPosition(X,Y)}] and
        [{towards((X1,Y1),(X,Y))}] and not [{obstacle(X1,Y1)}]
then [{stepTo(X1,Y1)}]
```

define communicateGold {

```
when [{at(X,Y)}] and not [{goldAt(X,Y)}] and [{loaded}]
then [{broadcast(goldAt(X,Y))}]
```

```
when not [{batteryLow}] then {
    movementToPosition /* 1 */
```

```
communicateGold /* 2 */
```

```
} else {
```

```
when not [{desiredPosition(X_R, Y_R)}] then {
[{adopt(desiredPosition(X_R, Y_R))}] /* 3 */
```

} else {

```
movementToPosition /* 4 */
```



Example: translation to the core language

 $\neg Q('batteryLow') \land Q('at(X,Y)') \land \neg Q('desiredPosition(X,Y)') \land Q('towards((X1,Y1),(X,Y))') \land \neg Q('obstacle(X1,Y1)')$ $\longrightarrow U('stepTo(X1,Y1)')$

 $\neg Q('batteryLow') \land Q('at(X,Y)') \land \neg Q('goldAt(X,Y)') \land Q('loaded')$ $\longrightarrow U('broadcast(goldAt(X,Y))')$

/* 3 */

/* 2 */

 $\neg \neg Q('batteryLow') \land \neg Q('desiredPosition(X_R, Y_R)') \\ \longrightarrow U('adopt(desiredPosition(X_R, Y_R))')$

 $/* 4 */ \neg \neg Q('batteryLow') \land \neg \neg Q('desiredPosition(X_R, Y_R)') \land Q('at(X,Y)') \land \neg Q('desiredPosition(X,Y)') \land Q('towards((X1,Y1),(X,Y))') \land \neg Q('obstacle(X1,Y1)') \rightarrow U('stepTo(X1,Y1)')$

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Related work

Modularity of rule-based agent programming languages:

3APL

- Dastani et. al.: Enacting and deacting roles in agent programming, 2004
- Riemsdijk et. al.: Goal oriented modularity in agent programming, 2006

AgentSpeak(L)

 Hübner et. al.: Programming declarative goals using plan patterns, 2006

GOAL

Hindriks: Modules as policies, ProMAS 2007



Ongoing & future-work

integration with BDI framework

- modular BDI architecture
- BDI rather a methodological guideline, than an explicit programming language?
- testing in real-world
 - Jazzyk language interpreter ~~> working prototype(!)
 - integration with Python, Answer Set Programming solver, later Prolog and LISP
 - demo application(!)



Summary

Structure of rule-based programming language:

- concept of mental state transformer
- powerful macros

Purely syntactical approach ~~> No change of the core language semantics!

Thank you for your attention.

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Summary

Structure of rule-based programming language:

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