

Code patterns for agent-oriented programming

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Motivation

- reactivity vs. deliberation ~→ hybrid architectures ~→ BDI
- programming with mental attitudes: beliefs, goals, etc.

agent oriented programming languages

- choose a set of agent-oriented features
 implement the set in the language interpreter
- fixed set of language constructs
 fixed architecture of created agent systems

extensions require changes of the language semantics \Rightarrow adaptation of the interpreter



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Problem & the way to go...

How to design extensible programming languages for cognitive agents. **2**

How to develop domain independent high level language constructs for programming with mental attitudes?



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Behavioural State Machines

A programming framework with clear separation between *knowledge representation* and agent's *behaviours*.

heterogeneous knowledge bases

structured source code, macros

```
the core \rightsquigarrow KR module \mathcal{M}
BSM agent system \rightsquigarrow \mathcal{A} = (\mathcal{M}_1, \dots, \mathcal{M}_n, \mathcal{P})
```

```
/* PICK an item behaviour */

when \models_{\mathcal{G}} [\{ task(pick(X)) \}] and <math display="inline">\models_{\mathcal{B}} [\{ see(X) \}] then {

when <math>\models_{\mathcal{B}} [\{ dir(X, Angle) \}] then \oslash_{\mathcal{E}} [\{ turn Angle \}] |

when <math>\models_{\mathcal{B}} [\{ dir(X, 'ahead'), dist(X,Dist) \}] then {

<math>\oslash_{\mathcal{E}} [\{ move forward Dist \}] \circ

\bigoplus_{\mathcal{B}} [\{ holds(X) \}]

}
```

/* either turn to the item, or */ /* pick up the item */

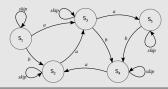
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BSM semantics

BSM ~~ labelled transition system



run of
$$\mathcal{P}$$
 $\lambda = \underbrace{s_1 \xrightarrow{a} s_3 \xrightarrow{b} s_4 \xrightarrow{\text{skip}} s_4 \xrightarrow{a} s_2}_{\mathcal{P}_1} \rightarrow \cdots$

reasoning about computation runs:

a logic interpreted over the same structure

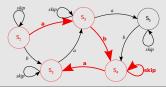
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BSM semantics BSM ~~ labelled transition system

■ operational ~→ computation runs



run of
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reasoning about computation runs:

→ a logic interpreted over the same structure!

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DCTL*=Dynamic Logic + CTL* $\theta ::= p | \neg \theta | \theta \land \theta | [\tau] \varphi$ $\varphi ::= \theta | \neg \varphi | \varphi \land \varphi | \bigcirc \varphi | \varphi \mathcal{U} \varphi | \varphi \mathcal{C} \varphi$ $[\tau] \varphi \rightsquigarrow during execution of <math>\tau, \varphi$ holds

From BSM to DCTL*: annotations \mathfrak{A} Annotated BSM $\mathcal{A}^{\mathfrak{A}} = (\mathcal{M}_1, \dots, \mathcal{M}_n, \mathcal{P}, \mathfrak{A})$ $\mathfrak{A} : \mathcal{Q}(\mathcal{A}) \cup \tau(\mathcal{A}) \rightarrow DCTL^*$

from subprograms to complex programs ~ aggregation

semantic characterization ---- the key to code re-usability

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Annotated BSM $\mathcal{A}^{\mathfrak{A}} = (\mathcal{M}_1, \dots, \mathcal{M}_n, \mathcal{P}, \mathfrak{A})$

 $\mathfrak{A}:\mathcal{Q}(\mathcal{A})\cup\tau(\mathcal{A})\to DCTL^*$

from subprograms to complex programs ---- aggregation

semantic characterization ~> the key to code re-usability

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Agent system architecture $\mathcal{A} = (\mathcal{B}, \mathcal{G}, \mathcal{E}, \mathcal{P})$



robot in a 3D environment: search & deliver

Structure:

- \mathcal{B} : belief base ($\models_{\mathcal{B}}, \oplus_{\mathcal{B}}, \ominus_{\mathcal{B}}$)
- \mathcal{G} : goal base ($\models_{\mathcal{G}}, \oplus_{\mathcal{G}}, \ominus_{\mathcal{G}}$)
- \mathcal{E} : interface to the environment \rightsquigarrow body ($\models_{\mathcal{E}}, \oslash_{\mathcal{E}}$)

Basic capabilities:

FIND: $[FIND]\mathfrak{A}(FIND) \Rightarrow [FIND^*] \diamond holds(item 42)$

 $\textbf{RUN_AWAY:} \ [\texttt{RUN_AWAY}] \mathfrak{A}(\texttt{RUN_AWAY}) \Rightarrow [\texttt{RUN_AWAY}^*] \Diamond \mathit{safe}$





BSM design patterns: TRIGGER

define TRIGGER($\varphi_{\mathbf{G}}, \tau$) when $\models_{\mathcal{G}} \varphi_{\mathbf{G}}$ then τ end

$\mathfrak{A}(\models_{\mathcal{G}}\varphi_{\mathbf{G}}) \rightarrow [\mathsf{TRIGGER}(\varphi_{\mathbf{G}},\tau)^*] \diamondsuit \mathfrak{A}(\tau)$

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BSM design patterns: ADOPT/DROP

$$\begin{array}{l} \operatorname{define} \mathsf{ADOPT}(\varphi_{\mathbf{G}},\psi_{\oplus}) \\ \quad \mathsf{when} \models_{\mathcal{B}} \psi_{\oplus} \text{ and not} \models_{\mathcal{G}} \varphi_{\mathbf{G}} \text{ then } \oplus_{\mathcal{G}} \varphi_{\mathbf{G}} \\ \operatorname{end} \end{array}$$

 $\begin{array}{l} \text{define DROP}(\varphi_{\mathbf{G}},\psi_{\ominus}) \\ \text{when } \models_{\mathcal{B}} \psi_{\ominus} \text{ and } \models_{\mathcal{G}} \varphi_{\mathbf{G}} \text{ then } \ominus_{\mathcal{G}} \varphi_{\mathbf{G}} \\ \text{end} \end{array}$



 $\begin{aligned} \mathfrak{A}(\models_{\mathcal{B}}\psi_{\oplus}) &\to [\mathsf{ADOPT}(\varphi_{\mathbf{G}},\psi_{\oplus})^*] \Diamond \mathfrak{A}(\models_{\mathcal{G}}\varphi_{\mathbf{G}}) \\ \mathfrak{A}(\models_{\mathcal{B}}\psi_{\ominus}) &\to [\mathsf{DROP}(\varphi_{\mathbf{G}},\psi_{\ominus})^*] \Diamond \neg \mathfrak{A}(\models_{\mathcal{G}}\varphi_{\mathbf{G}}) \end{aligned}$

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BSM design patterns: ACHIEVE

BG

 $\begin{array}{l} \text{define ACHIEVE}(\varphi_{\mathbf{G}},\varphi_{\mathbf{B}},\psi_{\oplus},\psi_{\ominus},\tau) \\ \text{TRIGGER}(\varphi_{\mathbf{G}},\tau) \mid \\ \text{ADOPT}(\varphi_{\mathbf{G}},\psi_{\oplus}) \mid \\ \text{DROP}(\varphi_{\mathbf{G}},\varphi_{\mathbf{B}}) \mid \\ \text{DROP}(\varphi_{\mathbf{G}},\psi_{\ominus}) \\ \text{end} \end{array}$

$[\mathsf{ACHIEVE}(\varphi_{\mathbf{G}},\varphi_{\mathbf{B}},\psi_{\oplus},\psi_{\ominus},\tau)^*]\mathfrak{A}(\models_{\mathcal{G}}\varphi_{\mathbf{G}})\mathcal{U}\mathfrak{A}(\models_{\mathcal{B}}\varphi_{\mathbf{B}}\lor\models_{\mathcal{B}}\psi_{\ominus})$

running example cont.

ACHIEVE(

```
achieve(has(item42)),
holds(item42),
needs(item42),
¬needs(item42) ∨ ¬exists(item42),
FIND)
```

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BSM design patterns: MAINTAIN

define MAINTAIN($\varphi_{\mathbf{G}}, \varphi_{\mathbf{B}}, \tau$) when not $\models_{\mathcal{B}} \varphi_{\mathbf{B}}$ then TRIGGER($\varphi_{\mathbf{G}}, \tau$) | ADOPT($\varphi_{\mathbf{G}}, \top$) end



 $\mathfrak{A}(\models_{\mathcal{G}}\varphi_{\mathbf{G}}) \to [\mathsf{MAINTAIN}(\varphi_{\mathbf{G}},\varphi_{\mathbf{B}}\tau)^*] \Box (\neg \mathfrak{A}(\models_{\mathcal{B}}\varphi_{\mathbf{B}}) \to \Diamond \mathfrak{A}(\models_{\mathcal{B}}\varphi_{\mathbf{B}}))$

running example cont.

MAINTAIN(*maintain*(*keep_safe*), *safe*, **RUN_AWAY**)

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Putting it altogether

Robot program

PERCEIVE o

MAINTAIN(

maintain(keep_safe),
threatened,
RUN_AWAY) |

ACHIEVE(

```
achieve(has(item42)),
holds(item42),
needs(item42),
¬needs(item42) ∨ ¬exists(item42),
FIND)
```



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Summary

different applications require different programming constructs



extensible agent oriented programming languages

purely syntactic approach to development of arbitrary high level programming constructs

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Thank you for your attention.

http://jazzyk.sourceforge.net/



see you at the poster session...

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